

# Hyperspeed Review of [B].

We inherit notations from KC's talk on [AB]. Today our goal is to construct monoidal equivalences

$$D(\check{I} \backslash \check{G}(k) / \check{I}) \cong \text{IndCoh}(\tilde{N} \times \tilde{N})$$

Notation:  $\tilde{g} \tilde{g} \tilde{g} =: st$ ,  $\tilde{N} \times \tilde{N} =: st^d$   
 $\tilde{g} \tilde{g} \tilde{N} =: st'$

i.e. generated

by pullback from  $D_{\text{mon}} \check{I} \check{I}$ .

$$\rightarrow D_{\text{mon}}(\check{I} \backslash \check{G}(k) / \check{I}) \cong \text{IndCoh}(\tilde{g} \times \tilde{g} / \tilde{G})$$

where we remind that  $\tilde{N}$  is the Springer resolution and  $\tilde{g}$  is the Grothendieck-Springer resolution.

$$D(\check{I} \backslash \check{G}(k) / \check{I}) \cong \text{IndCoh}(\tilde{g} \times \tilde{N} / \tilde{G})$$

In his senior thesis, James Tao demonstrated how to streamline a lot of [B] using higher category theory. Much of this talk will be copy-paste from his thesis.

## Recalling AB.

Recall that we got a functor  $\Phi_{\text{diag}}: \text{Coh}(G \backslash \tilde{N}) = \text{Coh}(B/\mathfrak{a}) \rightarrow \text{Dmod}(\check{I} \backslash \check{G}(k) / \check{I})$  via the following procedure:

- get  $\text{Coh}(pt/T)$  from Wakimoto.
  - get  $\text{Coh}(pt/G)$  via central sheaves.
  - ↓ upgrade to  $\text{Coh}(g/G)$  via log of monodromy
- } Prinfeld-Pluecker to get  $\text{Coh}(pt/B)$  (using Wakimoto filtration of central sheaves).

(We also get  $\Phi_{\text{IW}}: \text{Coh}(G \backslash \tilde{N}) \rightarrow \text{Dmod}(\check{F}\ell)(\check{I} \check{a}, X, \mathfrak{a})$  which intertwines the action on both sides. I'll be then combine. forgetting these very often...

The plan of attack is as follows:

- 1) Results in derived geometry yields that

$$\text{Perf}(St^d/G) \cong \text{Perf}(\tilde{N}/G) \otimes \text{Perf}(\tilde{N}/G) \otimes \text{Perf}(\tilde{g}/G)$$

So a map out of it can be obtained via a  $\tilde{g}/G$ -compatible ~~map~~ map out of  $\text{Perf}(\tilde{N}/G)$ .

In particular we obtain ~~the~~ by letting  $\bullet \tilde{N}/G$  acting on both sides via  $\Phi_{\text{diag}}$  on a tilting object  $\Xi$ .  
 a map to  $D(\check{F}\ell) \check{I}$

2) Given any functor  $\mathbb{F}: C \rightarrow D$  of DGcat, we can always form the ind-adjoint.

$$D \xrightarrow{\text{Yoneda}} \text{Funct}(D^{\text{op}}, \text{Vect}) \xrightarrow{\mathbb{F}^{\text{op}}} \text{Funct}(C^{\text{op}}, \text{Vect})$$

which gives a map  $\tilde{\mathbb{F}}: D(\tilde{\mathbb{F}}) \rightarrow \text{Incl}(\text{Perf}(St^d/G) \simeq \text{Qcoh}(St^d/G))$ .

~~that this factor is not the same as the one above~~

Claim: explicit cohomological vanishing can guarantee ~~that it's an iso~~ that it's an iso onto Coh.

2.5) One shows that conditions above can be rephrased purely in terms of  $\tilde{\mathbb{F}}$ .

3) To check the monoidal/module category claims we will need to invoke some compatibilities with  $\tilde{\mathbb{F}}$  to construct another functor  $\tilde{\mathbb{F}}': (\text{automorphic}) \rightarrow (\text{spectral})$  and show it coincides with  $\tilde{\mathbb{F}}$  above.

Remark on Free-monodromic Schemes

Recent development in higher algebra now allows us to ~~more~~ largely bypass the use of free-monodromic objects in this story. For completeness, however, we include a brief explanation of what's going on.

The difficulty lies in accessing  $\text{Coh}(\tilde{N}_g \times \tilde{N}_g)^G$ , which involves a derived tensor product.

To get around this issue, one can try to consider  $\tilde{N}_g \times \tilde{N}_g \xrightarrow{L} \hat{St} := \text{formal completion of } St \text{ over } \mathbb{D}_{N \in \mathfrak{g}}$ .

Claim:  $L^*$  IndCoh exists. (standard Koszul duality:  $\mathbb{D}^2$  pullback is equivalent to  $\mathbb{D}^2 K$ ; now  $K$  has a finite resolution).

Corollary:  $(L^* \text{IndCoh}, L_* \text{IndCoh})$  restricts to an adjoint pair  $(L^*, L_*)$  between Coh.

Since  $L_*$  is conservative, BBL says we can express  $\text{Coh}(\tilde{N}_g \times \tilde{N}_g)^G$  as  $(L_* L^*)$ -mod in  $\text{Coh}(\hat{St})^G$ ; the latter does not need derived geometry to express.

By some reason,  $L^*$  of  $\text{Coh}(\hat{St})^G$  generates the category, so if we can express  $\text{Coh}(\hat{St})^G$  on the dual side we'd be able to ~~obtain the desired iso~~ obtain the desired iso by matching monoids. Now since  $\hat{St}$  is an indscheme, its Coh contains some "pro-objects" (e.g. structure sheaf) and these turn out to be the pro-monodromic objects on the dual side.

More precisely, let  $\hat{D}$  be a certain pro-completion of  $\mathbb{D}(Y_0(G)/I_0)$  at the origin. (cf [B-YA-3-1])

~~I want to say~~ ~~but not exactly~~

then we have  $\hat{D} \cong \text{Indush}(\hat{S}t/G)$ , which is an upgrade of the statement regarding  $D_{\text{mon}}$ . ~~Also~~ This can be established by upgrading the entire  $[AB]$  to this setting.

\*: let  $X \rightarrow Y$  be a  $T$ -torsor, then we get the Atiyah bundle

$$0 \rightarrow \text{Lie}(T) \otimes \mathcal{O}_Y \rightarrow \text{At} \rightarrow \mathcal{O}_Y \rightarrow 0.$$

Let  $\mathcal{D}$  be the  $\mathcal{O}_Y$ -algebra generated by  $\text{At}$ . then  $D_{\text{mon}}(X) = \{ \mathcal{D}\text{-modules on which } \text{Lie}(T) \text{ acts nilpotently} \}$ .

Since  $\text{Lie}(T)$  is central, we can consider

$$\hat{\mathcal{D}} := \widehat{\text{Sym}(\text{Lie}(T))} \otimes \mathcal{D}, \text{ where } \wedge \text{ is completion w.r.t ideal } \text{Lie}(T).$$

and  $\hat{\mathcal{D}}$  above is roughly (how?) the category of  $\hat{\mathcal{D}}$ -modules.

The monodromy on  $\hat{\mathcal{D}}$  is then the action by  $\text{Lie}$  of monodromy, and one proceeds to match them. We shall not go down this path in this talk.

### Tilting Objects

In any highest weight category  $\mathcal{C}$ , an object  $T$  is tilting if it admits both a standard flag and a costandard flag.

Prop. (Tilting Exercises)

- ~~Tilting~~ Indecomposable tilting objects are in bijection with the standards.
- They generate  $\mathcal{C}$ .
- For any two tilting objects  $T_1, T_2$ ,  $\text{Ext}^n(T_1, T_2) = 0 \quad \forall n > 0$ .

$\Rightarrow K^b(T) = D^b(\mathcal{C})$ .  $\leftarrow$  This allows [B] to bypass difficulties with derived categories. We shall not worry about this aspect.

By this proposition, for any  $w \in W_{\text{aff}}$  we have a corresponding tilting object  $T_w \in D(\tilde{Fl})^i$ . Our main player is  $\Xi := T_{w_0}$ . It is also the projective cover of  $\Delta_0$ , known as the "big projective" in category  $\mathcal{O}$ .

Let  $Av_{Iw} : D(\check{F}\ell)^{\check{I}} \rightarrow D(\check{F}\ell)^{(I_0^{\vee}, X)}$

and  $Av_{Iw \rightarrow I\phi} : D(\check{F}\ell)^{(I_0^{\vee}, X)} \rightarrow D(\check{F}\ell)^{\check{I}}$

be the  $\check{\cdot}$ -averaging functors.

Prop. (Right convolution with  $\Xi$ ) =  $(Av_{Iw \rightarrow I_0} \circ Av_{I\phi \rightarrow Iw})$ .

(Compare in category  $\mathcal{O} : (- \circ \Xi) = (T_{\rho \rightarrow \lambda} \circ T_{\lambda \rightarrow -\rho})$ .)

(We might prove this if we have time.)  $\swarrow$  (left conv.) (see back side)

Construction of  $\mathbb{F}_{perf} : Perf(\mathcal{S}t^d/G) \rightarrow D(\check{F}\ell)^{\check{I}}$

A derived stack  $X$  is called perfect if  $\mathcal{O}_{\text{glob}}(X) = \text{Ind Perf}(X)$  ~~with affine diagonal~~

Equivalently, this means  $\mathcal{O}$  is compact and  $\mathcal{O}_{\text{glob}}$  is compactly generated.

e.g. - any derived qc scheme.

- (quasi-)projective scheme

$\Rightarrow \tilde{N}/G, g/G$  are perfect.

[BZFN] For  $X \rightarrow S, Y \rightarrow S$  maps of perfect stacks,  $Perf(X \times_S Y) \simeq Perf(X) \otimes_{Perf(S)} Perf(Y)$ .

Thus  $Perf(\mathcal{S}t^d/G) \simeq Perf(\tilde{N}/G) \otimes_{Perf(g/G)} Perf(\tilde{N}/G)$ .

So to get  $\mathbb{F}_{perf}$ , suffices to map from  $Perf(\tilde{N}/G)^2$  in a way compatible with  $Perf(g/G)$ -action. Namely, we let both sides map to  $D(\check{F}\ell)^{\check{I}}$  then cut on  $\Xi$  via convolution. Compatibility was explained in KC's talk.

Prop.  $\mathbb{F}_{perf}$  is fully faithful. ([13.6.2])

Pf sketch. Suffices to check on generators. Conv by Wakimoto / central sheaves have obvious inverses, so we reduce to check

$$\text{Hom}(\mathcal{O}_{\mathcal{S}t^d}(0, \mu), V^* \otimes \mathcal{O}_{\mathcal{S}t^d}(\lambda, 0)) \simeq \text{Hom}(\Xi * J_{\mu}, J_{\lambda} * * Z_{V^*} * \Xi)$$

Left Conv: Under the equivalence  $\Phi_{IW} : \text{Coh}(\tilde{N}/G) \xrightarrow{\sim} \text{Dmod}(\check{F}\ell)^{(\mathbb{Z}_0, \mathbb{Z}_1), \text{lc}}$   
 Convolution from left by  $\square$  corresponds to  $p^*p_*$ , where  $p: \tilde{N}/G \rightarrow g/G$  is the projection.

Now we use the conv. description above:

$$\begin{aligned} \text{RHS} &= \text{Hom}(A_{V_{I_0 \rightarrow IW}}(\Sigma * \bar{J}\mu), A_{V_{I_0 \rightarrow IW}}(\bar{J}\lambda * \bar{Z}V)) \quad (\text{we use the fact that } A_{V_{IW \rightarrow I_0}} \circ A_{V_{I_0 \rightarrow IW}} \\ &= \text{Hom}_{\text{Coh}(\tilde{N}/G)}(p^*p_* \mathcal{O}(\mu), \mathcal{O}(\lambda) \otimes V) \quad = A_{V_{IW \rightarrow I_0}}^* \circ A_{V_{I_0 \rightarrow IW}} \\ &= \text{Hom}_{\text{Coh}(\tilde{N}/G)}(p_{S,2}^* p_{S,1}^* \mathcal{O}(\mu), \mathcal{O}(\lambda) \otimes V) \quad (\text{since } \Phi_{IW} = A_{V_{I_0 \rightarrow IW}} \circ \Phi_{\text{diag}}) \\ &= \text{Hom}_{\text{Coh}}(p_{S,1}^* \mathcal{O}(\mu), p_2^*(\mathcal{O}(\lambda) \otimes V)) \quad \# \quad p_{S,i} : \text{St}^d/G \rightarrow \tilde{N}/G. \\ & \quad (p_{S,i}^* = p_{S,i}^{\dagger}) \\ & \quad \text{h/c } \omega = \mathcal{O}[d] \text{ in both source \& target.} \end{aligned}$$

### Construction and iso of $\Phi$

As explained in introduction, we always have the incl-right adjoint

$$\Phi : \text{Dmod}(\check{F}\ell)^{\mathbb{Z}_0, \mathbb{Z}_1} \rightarrow \text{Incl Perf} \cong \text{QCoh}(\text{St}^d/G).$$

The following allows us to say when it ~~is an iso to~~  $\text{Coh}$ .

Prop. ~~Let~~ Let  $\Phi : \text{Perf}(\mathbb{Z}/H) \rightarrow \mathcal{C}$  and  $\Phi : \mathcal{C} \rightarrow \text{Incl}(\text{Perf}(\mathbb{Z}/H)) \cong \text{QCoh}(\mathbb{Z}/H)$  be as above. Assume  $\Phi$  fully faithful and  $\mathcal{C}$  has bounded t-structure.

Further assume:

- 1)  $\Phi$  levels in  $\text{Coh}$ .
- 2)  $\Phi$  has finite cohomological amplitude.
- 3)  $\exists d > 0$  s.t.  $\forall M \in \mathcal{C}, H^i(\Phi(M)) = 0 \forall |i| \leq d \Rightarrow H^0(M) = 0$ .

Then  $\Phi$  factors through  $\mathcal{C} \cong \text{Coh}(\mathbb{Z}/H)$ .

Let's explain why this can be true. ~~for any  $F \in \text{Coh}(Z)$~~  for any  $F \in \text{Coh}(Z)$  we can find

$$(F^{\otimes N} \rightarrow F^{\otimes N-1} \rightarrow \dots \rightarrow F \rightarrow \dots)$$

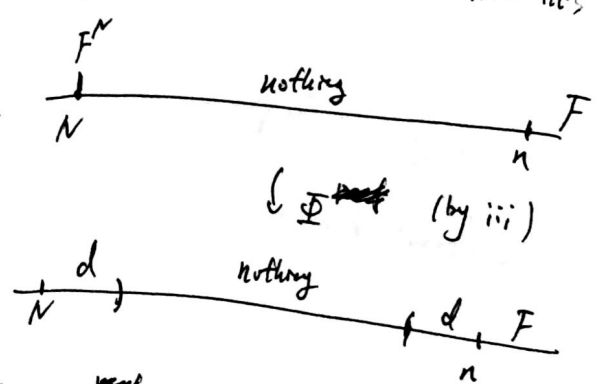
(take  $E$  locally free  $\rightarrow F$  aka Hilbert syzygy, then stupidly truncate)

where  $F^{\text{perf}}$  is in Perf and  $F^{\otimes N}$  is concentrated in deg  $N$ .

Since  $\mathbb{I}$  lands in Coh,  $\text{Hom}(\mathbb{I}(A), \mathbb{I}(B)) \simeq \text{Hom}(E, \mathbb{I}(B))$  for some good perf-approx.  $E$  of  $\mathbb{I}(A)$ .  $\mathbb{I}$  being fully faithful implies this  $\simeq \text{Hom}(\mathbb{I}(E), B)$ . At the same time,  $E \rightarrow \mathbb{I}(A) \hookrightarrow \mathbb{I}(E) \rightarrow A$ , which is a good approximation of  $A$  itself by iii), so actually  $\text{Hom}(A, B) = \text{Hom}(\mathbb{I}(E), B)$ . So fully faithful.

To get ess. surj., choose  $N \ll 0$ , then by coh. assumption,  $F^{\text{perf}}$  looks like this

$$F^{\otimes N} \rightarrow F^{\text{perf}} \rightarrow F \quad (F \in \text{Coh}^{\geq n})$$



But how do we use this?

The criterion above concerns cohomological vanishing in  $\text{Coh}(Z/H)$ , but how do we detect those?

This turns out to be fairly explicit if  $Z$  is projective over an affine base, because we can reduce to affine, for which the criterion is explicit: it amounts to computation of  $\text{Hom}(O, F[i])$ .

Since  $\text{obl}_Z: \text{Coh}(Z/H) \rightarrow \text{Coh}(Z)$  is t-exact, it suffices to compute

$$\text{Hom}_Z(O_Z, \text{obl}_Z F) \stackrel{\text{standard pull-push}}{=} \text{Hom}_{Z/H}(O_{Z/H}, F \otimes O_H)$$

This reduces us to the quasi-affine case. Writing  $Z \hookrightarrow \text{Proj } A$  and let  $\bar{Z} = \text{Spec } A/\mathfrak{m}$ . We have an open embedding  $j: Z \hookrightarrow \bar{Z}$ . The following is a derived version of Serre vanishing theorem:

- Prop 1) For  $F \in \text{Coh}(Z)$ ,
- for  $k \gg 0$ ,  $\forall d$ ,  $H^d_{g_{\geq k}} j_* F = 0 \Leftrightarrow H^d F = 0$ .
  - 2) for any  $k$ ,  $\forall F \in \text{Coh}(Z)$ ,  $F \in \text{Coh}(Z) \hookrightarrow g_{\geq k} j_* F \in \text{Coh}(\bar{Z})$ .

Then relevant identity vanishing is given by considering

$$\Gamma(\text{Spec } A, \mathcal{O}_{\mathbb{P}^k} \otimes \mathcal{I}(n))$$

$$= \bigoplus_{n \geq k} \Gamma(\mathbb{P}^k, (\mathcal{O}_{\mathbb{P}^k} \otimes \mathcal{I}(n)) \otimes \mathcal{O}_{\mathbb{P}^k}(n))$$

$$= \bigoplus_{n \geq k} \text{Hom}(\mathcal{O}_{\mathbb{P}^k}(-n), \mathcal{I}(n) \otimes \mathcal{O}_{\mathbb{P}^k}(n))$$

$$= \bigoplus_{\substack{n \geq k \\ \lambda \in \Lambda^+}} \text{Hom}(\mathcal{O}_{\mathbb{P}^k}(-n) \otimes V_\lambda, \mathcal{I}(n) \otimes V_\lambda) = \bigoplus_{\substack{n \geq k \\ \lambda \in \Lambda^+}} \text{Hom}_{D(\check{F}_1)^{\check{i}, \text{lc}}}(\check{\mathcal{P}}_{\text{perf}}(\mathcal{O}_{\mathbb{P}^k}(-n) \otimes V_\lambda), \mathcal{I}(n) \otimes V_\lambda)$$

from which ~~it~~ now follows from our knowledge about  $\check{\mathcal{P}}_{\text{diag}}$ .

Putting everything above together, we have obtained an equivalence  $D(\check{F}_1)^{\check{i}} \cong \text{IndCoh}(\text{St}^d/G)$  as plain DG categories. Let  $\check{\mathcal{F}}: \text{Coh}(\text{St}^d/G) \xrightarrow{\sim} D(\check{F}_1)^{\check{i}, \text{lc}}$  be the opposite direction.

Recovering the monoidal structure

It turns out there's a much faster way of obtaining  $\check{\mathcal{F}}$ . Namely, one can prove w/o much difficulty that  $\text{Perf}(\tilde{N}/G)$  is ~~also~~ self-dual as  $\text{Perf}(\mathfrak{g}/G)$ -module, and we get (using  $\tilde{N} \rightarrow \mathfrak{g}$  proper scheme),  $\text{Coh}(\text{St}^d/G) \cong \text{Perf}(\tilde{N}/G) \otimes_{\text{Perf}(\mathfrak{g}/G)} \text{Perf}(\tilde{N}/G) \cong \text{Hom}_{\text{Perf}(\mathfrak{g}/G)}(\text{Perf}(\tilde{N}/G), \text{Perf}(\tilde{N}/G))$  which is moreover a monoidal equivalence.

Now observe that  $\check{\mathcal{F}}_{\text{inv}}: \text{Perf}(\tilde{N}/G) \cong \text{Coh}(\tilde{N}/G) \xrightarrow{\check{\mathcal{F}}_{\text{diag}}} D(\check{F}_1)^{\check{i}} \xrightarrow{A_{\check{I}_0 \rightarrow \check{I}_1}} D(\check{F}_1)^{\check{i}, X}$  can be upgraded to a map of  $\text{Perf}(\mathfrak{g}/G)$ -module categories. At the same time, action of  $D(\check{F}_1)^{\check{i}}$  on  $D(\check{F}_1)^{\check{i}, X}$  ~~is~~ commutes with  $\text{Perf}(\mathfrak{g}/G)$  action via central elements.

Thus the action map provides another functor  $\check{\mathcal{F}}': D(\check{F}_1)^{\check{i}, \text{lc}} \rightarrow \text{Coh}(\text{St}^d/G)$  (It's not clear how to show directly this is an equivalence.) Our job now is to show that  $\check{\mathcal{F}}' \cong \check{\mathcal{F}}$  as plain functors.



Recall the inclusion of a full subcategory  $C_0 \hookrightarrow C$  is called dense if  $LKE_i = Id_C$ .  
 Typical example of this is the Yoneda embedding.

Prop. For  $X$  as discussed above (having the approximation property),  $Perf \hookrightarrow Coh$  is dense.  
 (This is a repackaging of the approximation property).

Then consider

$$\begin{array}{ccccc} Perf(S^d/G) & \hookrightarrow & Coh(S^d/G) & \xrightarrow{\sim} & Eucl(Perf(\tilde{N}/G)) \\ \downarrow i & & & \nearrow \text{act} & \\ Coh(S^d/G) & & & \xrightarrow{\sim} & Eucl(Perf(\tilde{N}/G)) \end{array}$$

we see  $act$  is LKEed from  $(act \circ i)$ .

Similarly,  $\Phi : Coh(S^d/G) \rightarrow D(\check{F}\ell)^{\check{I}}$  is LKEed from  $(\Phi \circ i) \simeq \Phi_{perf}$ .

Let  $D(\check{F}\ell)^{\check{I}} \rightarrow Eucl(-)$  be denoted  $act_D$ . Recall we want  $\Phi \circ \check{\Psi}' \simeq Id$ ,  
 i.e.  $act \simeq act_D \circ \Phi$ . By argument above, we get a natural transformation from

the following proposition:

Prop:  $act \circ i \simeq act_D \circ \Phi_{perf}$ . (This is proven using similar argument as fully faithfulness of  $\Phi_{perf}$ ).

Then to check the Nat is an iso, it suffices to do so on objects.

Unpacking definition we see that it amounts to check  $\forall F \in Coh(S^d/G), P \in Perf(\tilde{N}/G)$   
 we have:  $\Phi_{Iw}(F \star P) \simeq \Phi(F) \star \Phi_{Iw}(P)$ .

Noting that Nat between finite-amplitude functors is iso iff their restriction to  $Perf$  is,  
 it suffices to check that  $\Phi_{Iw}(- \star P)$  and  $\Phi(-) \star \Phi_{Iw}(P)$  have finite amplitude.  
 This follows from our knowledge about  $\Phi$  and  $\Phi_{Iw}$ .